PROVE IT YOURSELF THE BAKER-CAMPBELL-HAUSDORFF FORMULA

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ABSTRACT. Exercise sheets distributed at the summer school "Algebra, Topology and Fjords", Nord-fjordeid June 3-11, 2011.

1. NOTATION AND SET-UP

Let R be a unitary associative \mathbb{Q} -algebra and $I \subset R$ a nilpotent ideal. Denote by $1 + I = \{1 + a \mid a \in I\} \subset R$; notice that every element of 1 + I is invertible and $(1 + a)^{-1} = 1 + (-a + a^2 + \cdots) \in 1 + I$. Denote by $\operatorname{End}(R)$ the associative \mathbb{Q} -algebra of endomorphisms of R, considered as a vector space over \mathbb{Q} .

Define the **exponential**

$$e\colon I \to 1+I \subset R, \quad e^a = \sum_{n \ge 0} \frac{a^n}{n!},$$

and the ${\bf logarithm}$

log:
$$1 + I \to I$$
, $\log(1 + a) = \sum_{n=1}^{\infty} (-1)^{n-1} \frac{a^n}{n}$.

We assume already proved that exponential and logarithm are one the inverse of the other, i.e. for every $a, b \in I$ we have

$$\log(e^a) = a, \qquad e^{\log(1+b)} = 1+b.$$

2. The Hausdorff formula

Given $a \in R$ denote

$$a \colon R \to R,$$
 $(\operatorname{ad} a)(x) = [a, x] = ax - xa$

Exercise 2.1. Prove that for every $a, b \in R$ and $n \ge 0$ we have

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$$(\operatorname{ad} a)^n b = \sum_{i=0}^n (-1)^i \binom{n}{i} a^{n-i} b a^i = \sum_{i=0}^n \binom{n}{i} a^{n-i} b (-a)^i.$$

Deduce that if $a \in I$ then $\operatorname{ad} a$ is nilpotent in $\operatorname{End}(R)$ and therefore thre are defined the invertible operators

$$e^{\operatorname{ad} a} = \sum_{n \ge 0} \frac{(\operatorname{ad} a)^n}{n!} \in \operatorname{End}(R),$$
$$\frac{e^{\operatorname{ad} a} - I}{\operatorname{ad} a} = \sum_{n \ge 0} \frac{(\operatorname{ad} a)^n}{(n+1)!} \in \operatorname{End}(R),$$
$$\frac{\operatorname{ad} a}{e^{\operatorname{ad} a} - I} = \left(\frac{e^{\operatorname{ad} a} - I}{\operatorname{ad} a}\right)^{-1} = \sum_{n \ge 0} \frac{B_n}{n!} (\operatorname{ad} a)^n \in \operatorname{End}(R).$$

where B_n are the Bernoulli numbers.

Exercise 2.2. In the notation above prove that:

(1) For every $a \in I$ and $b \in R$

$$e^{\operatorname{ad} a}b := \sum_{n \ge 0} \frac{(\operatorname{ad} a)^n}{n!}b = e^a b e^{-a}.$$

- (2) For every $a \in I$ and $b \in R$ we have ab = ba if and only if $e^a b = be^a$.
- (3) For every $a, b \in I$ we have $e^a b = be^a$ if and only if $e^a e^b = e^b e^a$.
- (4) Given $a, b \in I$ such that ab = ba, then

$$e^{a+b} = e^a e^b = e^b e^a$$
, $\log((1+a)(1+b)) = \log(1+a) + \log(1+b)$.

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Let t be a variable and denote by $d: R[t] \to R[t], d(a) = a'$, the derivation operator:

$$(\sum a_n t^n)' = \sum na_n t^{n-1}$$

Multiplication on the left give an injective morphism of Q-algebras

$$\phi \colon R[t] \to \operatorname{End}(R[t]), \qquad \phi(a)b = ab.$$

Prove that:

$$\phi(a') = [d, \phi(a)], \qquad \forall \ a \in R[t]$$

 $\phi(a) = [a, \phi(a)], \qquad \forall \ a \in I[t], \\ \phi(e^{a}) = e^{\phi(a)}, \qquad \phi((e^{a})') = de^{\phi(a)} - e^{\phi(a)} d, \qquad \phi((e^{a})'e^{-a}) = d - e^{\operatorname{ad}\phi(a)} d, \qquad \forall \ a \in I[t], \\ \phi(e^{a}) = e^{\phi(a)} - e^{\phi(a)} d, \qquad \phi((e^{a})'e^{-a}) = d - e^{\operatorname{ad}\phi(a)} d, \qquad \forall \ a \in I[t], \\ \phi(e^{a}) = e^{\phi(a)} - e^{\phi(a)} d, \qquad \phi((e^{a})'e^{-a}) = d - e^{\operatorname{ad}\phi(a)} d, \qquad \forall \ a \in I[t], \\ \phi(e^{a}) = e^{\phi(a)} - e^{\phi(a)} d, \qquad \phi((e^{a})'e^{-a}) = d - e^{\operatorname{ad}\phi(a)} d, \qquad \forall \ a \in I[t], \\ \phi(e^{a}) = e^{\phi(a)} - e^{\phi(a)} d, \qquad \phi((e^{a})'e^{-a}) = d - e^{\operatorname{ad}\phi(a)} d, \qquad \forall \ a \in I[t], \\ \phi(e^{a}) = e^{\phi(a)} - e^{\phi(a)} d, \qquad \phi((e^{a})'e^{-a}) = d - e^{\operatorname{ad}\phi(a)} d, \qquad \forall \ a \in I[t], \\ \phi(e^{a}) = e^{\phi(a)} - e^{\phi(a)} d, \qquad \phi((e^{a})'e^{-a}) = d - e^{\operatorname{ad}\phi(a)} d, \qquad \forall \ a \in I[t], \\ \phi(e^{a}) = e^{\phi(a)} - e^{\phi(a)} d, \qquad \phi(e^{a}) = e^{\phi(a)} - e^{\operatorname{ad}\phi(a)} d, \qquad \forall \ a \in I[t], \\ \phi(e^{a}) = e^{\phi(a)} - e^{\phi(a)} d, \qquad \phi(e^{a}) = e^{\operatorname{ad}\phi(a)} d, \qquad \forall \ a \in I[t], \\ \phi(e^{a}) = e^{\operatorname{ad}\phi(a)} - e^{\operatorname{ad}\phi(a)} d, \qquad \forall \ a \in I[t], \\ \phi(e^{a}) = e^{\operatorname{ad}\phi(a)} d, \qquad \forall \ a \in I[t], \\ \phi(e^{a}) = e^{\operatorname{ad}\phi(a)} d, \qquad \forall \ a \in I[t], \\ \phi(e^{a}) = e^{\operatorname{ad}\phi(a)} d, \qquad \forall \ a \in I[t], \\ \phi(e^{a}) = e^{\operatorname{ad}\phi(a)} d, \qquad \forall \ a \in I[t], \\ \phi(e^{a}) = e^{\operatorname{ad}\phi(a)} d, \qquad \forall \ a \in I[t], \\ \phi(e^{a}) = e^{\operatorname{ad}\phi(a)} d, \qquad \forall \ a \in I[t], \\ \phi(e^{a}) = e^{\operatorname{ad}\phi(a)} d, \qquad \forall \ a \in I[t], \\ \phi(e^{a}) = e^{\operatorname{ad}\phi(a)} d, \qquad \forall \ a \in I[t], \\ \phi(e^{a}) = e^{\operatorname{ad}\phi(a)} d, \qquad \forall \ a \in I[t], \\ \phi(e^{a}) = e^{\operatorname{ad}\phi(a)} d, \qquad \forall \ a \in I[t], \\ \phi(e^{a}) = e^{\operatorname{ad}\phi(a)} d, \qquad \forall \ a \in I[t], \\ \phi(e^{a}) = e^{\operatorname{ad}\phi(a)} d, \qquad \forall \ a \in I[t], \\ \phi(e^{a}) = e^{\operatorname{ad}\phi(a)} d, \qquad \forall \ a \in I[t], \\ \phi(e^{a}) = e^{\operatorname{ad}\phi(a)} d, \qquad \forall \ a \in I[t], \\ \phi(e^{a}) = e^{\operatorname{ad}\phi(a)} d, \qquad \forall \ a \in I[t], \\ \phi(e^{a}) = e^{\operatorname{ad}\phi(a)} d, \qquad \forall \ a \in I[t], \\ \phi(e^{a}) = e^{\operatorname{ad}\phi(a)} d, \qquad \forall \ a \in I[t], \\ \phi(e^{a}) = e^{\operatorname{ad}\phi(a)} d, \qquad \forall \ a \in I[t], \\ \phi(e^{a}) = e^{\operatorname{ad}\phi(a)} d, \qquad \forall \ a \in I[t], \\ \phi(e^{a}) = e^{\operatorname{ad}\phi(a)} d, \qquad \forall \ a \in I[t], \\ \phi(e^{a}) = e^{\operatorname{ad}\phi(a)} d, \qquad \forall \ a \in I[t], \\ \phi(e^{a}) = e^{\operatorname{ad}\phi(a)} d, \qquad \forall \ a \in I[t], \\ \phi(e^{a}) = e^{\operatorname{ad}\phi(a)} d, \qquad \forall \ a \in I[t], \\ \phi(e^{a}) = e^{\operatorname{ad}\phi($

and deduce from the injectivity of ϕ that

$$(e^{a})'e^{-a} = \frac{e^{\operatorname{ad} a} - 1}{\operatorname{ad} a}(a').$$

Now, let $a, b \in I$ and define

$$Z = \log(e^{ta}e^b) \in I[t].$$

Prove that

$$\frac{e^{\operatorname{ad} Z} - 1}{\operatorname{ad} Z}(Z') = (e^Z)'e^{-Z} = a.$$

and therefore

Therefore Z = Z(t) is the solution of the Cauchy problem

$$Z' = \frac{\operatorname{ad} Z}{e^{\operatorname{ad} Z} - 1}(a) = \sum_{n \ge 0} \frac{B_n}{n!} (\operatorname{ad} Z)^n(a), \qquad Z(0) = Z_0 = b$$

where the B_n 's are the Bernoulli numbers $\left(\sum \frac{B_n}{n!} t^n = \frac{t}{e^t - 1}\right)$.

Theorem 2.3. Given $a, b \in I$ we have

$$e^{a}e^{b} = e^{a \bullet b}, \quad where \quad a \bullet b = \sum_{n \ge 0} Z_n,$$

and

$$Z_0 = b, \qquad Z_{r+1} = \frac{1}{r+1} \sum_{m \ge 0} \frac{B_m}{m!} \sum_{i_1 + \dots + i_m = r} (\operatorname{ad} Z_{i_1}) (\operatorname{ad} Z_{i_2}) \cdots (\operatorname{ad} Z_{i_m}) a.$$

Proof. Exercise. Hint $\log(e^{ta}e^b) = Z = Z_0 + tZ_1 + \dots + t^n Z_n + \dots$

The first terms of the above series are

$$a \bullet b = a + b + \frac{1}{2}[a, b] + \frac{1}{12}[a, [a, b]] + \frac{1}{12}[b, [a, b]] + \cdots$$

Since $(e^a e^b)e^c = e^a(e^b e^c)$ the product $I \times I \xrightarrow{\bullet} I$ is associative. If L is a Lie subalgebra of I and $a, b \in L$, then $a \bullet b \in L$ and $a \bullet b - a - b$ belongs to the Lie ideal generated by [a, b].

The formula of Theorem 2.3 allows to define for every nilpotent Lie algebra L a map

$$L \times L \to L, \qquad (a,b) \mapsto a \bullet b$$

commuting with morphisms of Lie algebras. Notice that if [a,b] = 0 then $a \bullet b = a + b$ and then $a \bullet (-a) = 0.$

3. Tree summation formula for BCH product

Recall that a tree is called a *rooted tree* if one vertex has been designated the *root*. Every rooted tree has a natural structure of directed tree such that, for every vertex u, there exists a unique directed path from u to the root. We shall write $u \to v$ if the vertex v belongs to the directed path from u to the root. A leaf is a vertex without incoming edges. A vertex is called *internal* if it is not a leaf.

From now on, we consider only planar binary rooted trees; we denote by \mathcal{B} the set of finite planar binary rooted trees with the root at the top and the leaves at the bottom (i.e., every directed path moves upward); binary means that every internal vertex has exactly two incoming edges.

We also write

$$\mathcal{B} = \bigcup_{n>0} \mathcal{B}_n,$$

where \mathcal{B}_n is the set of planar binary rooted trees with n leaves and, for every $\Gamma \in \mathcal{B}$, we denote by $L(\Gamma)$ the set of leaves of Γ . The planarity of the tree gives, for every internal vertex v, a total ordering of the edges ending on v, from the leftmost to the rightmost.

Definition 3.1. A rightmost branch of a planar binary rooted tree $\Gamma \in \mathcal{B}$ is a maximal connected subgraph $\Omega \subset \Gamma$, with at least two vertices and with the property that every edge of Ω is a rightmost edge of Γ .

Definition 3.2. A local rightmost leaf is a leaf lying on a rightmost branch. Given an internal vertex v, we call m(v) the leaf lying on the rightmost branch containing v. We also denote by d(v) the distance between v and m(v).

Definition 3.3. A subroot is the vertex of a rightmost branch which is nearest to the root. The set of subroots of a finite planar rooted tree Γ will be denoted by $R(\Gamma)$.

Let R be a (non associative) algebra over \mathbb{Q} and $\Gamma \in \mathcal{B}$. Labelling the leaves of Γ with elements of R, we can associate the product element in R obtained by the usual operadic rules, i.e., we perform the product of R at every internal vertex in the order arising from the planar structure of the directed tree. Given any map $f: L(\Gamma) \to R$ (the labelling), we denote by $Z_{\Gamma}(f) \in R$ the corresponding product element.

Definition 3.4. Given two leaves l_1 and l_2 in $\Gamma \in \mathcal{B}$, we say $l_1 \leq l_2$ if $l_1 = l_2$ or there exists a subroot $v \in R(\Gamma)$ such that $l_2 = m(v)$ and $l_1 \rightarrow v$.

Definition 3.5. For every poset (A, \leq) , we denote

$$\mathcal{B}(A) = \{ (\Gamma, f) \mid \Gamma \in \mathcal{B}, f : (L(\Gamma), \preceq) \to (A, \leq), f \text{ monotone } \}$$

In a similar way we define $\mathcal{B}_n(A)$, for every n > 0.

Definition 3.6. Let $b_n = B_n/n!$, or equivalently $\sum b_n t^n = \frac{t}{e^t - 1}$. Given a poset A and $(\Gamma, f) \in \mathcal{B}(A)$, let us define

$$c_{(\Gamma,f)} := \prod_{v \in R(\Gamma)} \frac{b_{d(v)}}{t(v)} ,$$

where for every subroot $v \in R(\Gamma)$, we have

t(v) = number of leaves $u \in L(\Gamma)$ such that $u \to v$ and f(u) = f(m(v)).

Theorem 3.7. Let L be a nilpotent Lie algebra, then for every $a, b \in L$ we have

(1)
$$a \bullet b = \sum_{(\Gamma, f) \in \mathcal{B}(b \le a)} c_{(\Gamma, f)} Z_{\Gamma}(f).$$

Proof. Left as exercise. Hint: let $\mathcal{A} \subset \mathcal{B}(b \leq a)$ be the subset of trees having every local rightmost leave labelled by a; then we have

$$\sum_{(\Gamma,f)\in\mathcal{B}(b\leq a)}c_{(\Gamma,f)}Z_{\Gamma}(f)=b+\sum_{(\Gamma,f)\in\mathcal{A}}c_{(\Gamma,f)}Z_{\Gamma}(f).$$

Now use Theorem 2.3 and induction on the number of subroots.

References: details about Hausdorff formula and its proof can be found in the book of B.C. Hall: *Lie Groups, Lie Algebras, and representations. An elementary introduction.* A detailed proof of Theorem 3.7 will appear in a forthcoming paper by D. Iacono and M. Manetti.