



# Cyclic colliding permutations

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## Abstract

We study lower and upper bounds for the maximum size of a set of pairwise cyclic colliding permutations.

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## 1 Preliminaries

We say that two permutations  $x, y \in S_n$  are *cyclic colliding* if and only if there exists an index  $1 \leq i \leq n$  such that the images of  $i$  according to  $x$  and  $y$  differ by 1 modulo  $n$ .

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More generally we consider

$$T_m(n) = \max\{|C| : C \subseteq S_n, \forall \{x, y\} \in \binom{C}{2} \exists i \in [n] : |x_i - y_i| \equiv 1 \pmod{m}\}.$$

We want to determine  $T_m(n)$  and  $T^*(n) = T_n(n)$  at least asymptotically. This is in analogy with a similar problem introduced by Körner and the second author in [3]: two permutations  $x, y \in S_n$  are *colliding* if and only if there exists an index  $1 \leq i \leq n$  such that the images of  $i$  by  $x$  and  $y$  differ by 1. The best known lower bound for  $T(n) := T_{n+1}(n)$ , that is the maximum size of a set of pairwise colliding permutations, can be found in [1].

We say that two permutations in  $T_m(n)$  are *m-colliding*. This encompasses both the definitions of cyclic colliding permutations (when  $m = n$ ) and colliding permutations (when  $m > n$ ).

The following is obvious.

**Proposition 1.1** *If  $m'$  divides  $m$ , then  $T_{m'}(n) \geq T_m(n)$ .*

We define the *parity pattern* ( $pp$ ) of a permutation  $x = (x_1, x_2, \dots, x_n)$  by  $pp(x) = (x_1[2], x_2[2], \dots, x_n[2])$ . For example, if  $x = \mathbf{1} := (1, 2, \dots, n)$  is the identical permutation, then  $pp(x) = (1, 0, 1, 0, \dots)$ . Since the parity pattern of a permutation is balanced binary sequence, there are  $\binom{n}{\lfloor n/2 \rfloor}$  possible parity patterns.

Setting  $xRy$  if and only if  $pp(x) = pp(y)$  defines an equivalence relation, with each class associated to a parity pattern. Clearly, if  $m$  is even, two  $m$ -colliding permutations cannot have the same parity pattern, i.e.  $x \rightarrow pp(x)$  restricted to a  $m = 2m'$ -colliding code is injective. Thus

**Proposition 1.2**  $T_{2m'}(n) \leq \binom{n}{\lfloor n/2 \rfloor}$ .

When  $m' = 1$ , equality holds, since two permutations belonging to different classes will be 2-colliding:

**Proposition 1.3**  $T_2(n) = \binom{n}{\lfloor n/2 \rfloor}$ .

We now focus on the case of cyclic collision ( $m = n$ ).

## 2 Lower bounds

It is immediate to see that if two permutations are colliding, then they are cyclic colliding, that is,  $T^*(n) \geq T(n)$ . This can be improved:

**Proposition 2.1**  $T^*(n) \geq 2T(n-1)$ .

**Proof.** Let  $C \subseteq S_{2,\dots,n}$  and  $D \subseteq S_{1,\dots,n-1}$  two sets of permutations of length  $n - 1$  pairwise colliding of maximum cardinality  $T(n - 1)$ . The set of permutations of  $[n]$

$$E := 1 \cdot C \cup n \cdot D$$

obtained by prefixing every permutation in  $C$  by 1 and every permutation in  $D$  by  $n$  is clearly pairwise cyclic colliding, and  $|E| = 2T(n - 1)$ .  $\square$

### 3 Upper bounds

We distinguish two cases, depending on the parity of  $n$ ; the even case  $n = m = 2m'$  follows from Proposition 1.2.

**Proposition 3.1**  $T^*(2m) \leq \binom{2m}{m}$ .

Now we analyze the case when  $n$  is odd.

**Proposition 3.2**  $T^*(2m + 1) \leq 3\binom{2m+1}{m}$ .

**Proof.** Let  $\sigma$  be a permutation of  $S_n, n = 2m + 1$ . The Hamming weight (i.e. the number of 1's) of  $pp(\sigma)$  is  $m + 1$ , thus there are  $\binom{2m+1}{m}$  parity patterns. Let now  $C$  be a cyclic colliding code; observe that in this case the map  $pp$  is no longer injective as in the case of  $n$  even: however we want to prove that  $pp$  is at most 3-to-1 when restricted on  $C$ . Without loss of generality, let  $z := (1, 1, \dots, 1, 0, 0, \dots, 0)$  be the parity pattern of some codeword, say:  $c^1 = (1, 3, 5, \dots, 2m + 1, 2, 4, \dots, 2m)$ . Let  $D := pp^{-1}(z) = \{c^1, c^2, \dots\}$  be the pre-image of  $z$  in  $C$ : we want to show that  $|D| \leq 3$ . Obviously the property of being cyclic colliding is inherited to subsets of a any cyclic colliding code (it is a pairwise condition holding for all pairs of the code); hence for  $D$  to be cyclic colliding, we must have: for  $i \neq j, c^i$  and  $c^j$  have the pair  $\{1, 2m + 1\}$  in some position (it is indeed the only way to be cyclic colliding and have the same parity pattern). Thus they never have a 1 nor a  $2m + 1$  in the same position.

Thus, without loss of generality, either:

$$\begin{aligned} c^1 &= ( \quad 1, \quad 2m + 1, \quad *, \quad *, \quad \dots, \quad *) \\ c^2 &= ( \quad 2m + 1, \quad 1, \quad *, \quad *, \quad \dots, \quad *) \end{aligned}$$

and  $|D| = 2$ ; or

$$\begin{aligned} c^1 &= ( \quad 1, \quad *, \quad 2m + 1, \quad *, \quad *, \quad \dots, \quad *) \\ c^2 &= ( \quad 2m + 1, \quad 1, \quad *, \quad *, \quad *, \quad \dots, \quad *) \\ c^3 &= ( \quad *, \quad 2m + 1, \quad 1, \quad *, \quad *, \quad \dots, \quad *) \end{aligned}$$

and  $|D| = 3$ . □

### Proposition 3.3

$$T^*(n) \leq nT(n-1).$$

**Proof.** We partition the permutations of a “code”  $C$  (that is, a family of permutations of  $n$  with the maximum cardinality with respect to the property of being pairwise cyclic colliding), according to the positions of the digit 1: let  $C_j = \{x = (x_1, \dots, x_n) \in C : x_j = 1\}$ , so that  $C = C_1 \cup \dots \cup C_n$ , with the  $C_j$ 's all disjoint (possibly empty). Each  $C_j$  contains permutations that are pairwise colliding, where the digits  $\{2, \dots, n\}$  are responsible for the collisions in  $C_j$  (since the cyclic collisions due to the digits 1 and  $n$  cannot appear in the  $C_j$ ), so that  $|C_j| \leq T(n-1)$ . □

### Corollary 3.4

$$2T(n-1) \leq T^*(n) \leq nT(n-1).$$

### Remarks and questions

- (i) In the case of cyclic collision, there is no proof of supermultiplicativity as for  $T(n)$ , namely :  $T(n+m) \geq T(n)T(m)$ ; thus the determination of  $T^*(n)$  cannot be seen as a “capacity” problem [1,3]. Setting  $R_n := (1/n) \log_2 T_n$ , we have by Fekete’s lemma that  $R_n$  tends to a limit  $R$  as  $n$  goes to infinity. Thanks to the previous corollary, we get directly the convergency of the analogous quantity in the cyclic case; furthermore,  $R^* = R$  holds.
- (ii) Can we prove/disprove that  $T^*(n) \leq T^*(n+1)$ ?
- (iii) Can we prove/disprove that  $T^*(n) \leq T(n+1)$ ?

We know the values of  $T(n)$  up to  $n = 9$  (the cases of 8 and 9 were found independently by Adolfo Piperno and Brik [2], communicated by Adriano Garsia), and they are both of the form  $\binom{n}{n/2}$ . For  $n = 10$ , A. Garsia and E. Sergel found through computer search different sets of pairwise colliding permutations consisting of 251 elements (one less than the upper bound  $\binom{10}{5} = 252$ ).

We found a code  $E \subseteq S_5$  of 20 pairwise cyclic colliding permutations of 5 elements, which improves on the lower bound of 12 given by Proposition 2.1. This construction is structured as follows. Let

$$E' = \{x = (x_1, \dots, x_5) : x \text{ is a cyclic shift of } (1, 3, 2, *, *)\},$$

that is

$$E' = \{(1, 3, 2, *, *), (*, 1, 3, 2, *), (*, *, 1, 3, 2), (2, *, *, 1, 3), (3, 2, *, *, 1)\}.$$

$E'$  consists of pairwise colliding permutations (as shown in Lemma 4.6 of [3]), hence cyclic colliding. One can “double” each partial permutation of  $E'$  filling the joker symbols  $*$  of  $E'$  once with 4, 5, then with 5, 4 in the order: call the corresponding sets of 5 permutations  $E'(4, 5)$  and  $E'(5, 4)$  respectively: putting them together, one obtains 10 pairwise colliding permutations (hence cyclic colliding). In a similar way, we build

$$E'' = \{x = (x_1, \dots, x_5) : x \text{ is a cyclic shift of } (5, 3, 4, *, *)\},$$

and “double” each of its elements filling the joker symbols  $*$  of  $E'$  once with 1, 2, then with 2, 1 in this order, to obtain  $E''(1, 2)$ ,  $E''(2, 1)$ , whose union leads to 10 colliding permutations. While the resulting set

$$F = E'(4, 5) \cup E'(5, 4) \cup E''(1, 2) \cup E''(2, 1)$$

is not a colliding code, it is surprisingly a cyclic colliding code.

We summarize the known values (or bounds) of the different considered types of  $T$ 's up to 10 in the following table.

$n$	1	2	3	4	5	6	7	8	9	10
$T_2(n) = \binom{n}{\lfloor n/2 \rfloor}$	1	2	3	6	10	20	35	70	126	252
$T(n)$	1	2	3	6	10	20	35	70	126	$251 \leq ? \leq 252$
$T^*(n)$	1	2	6	6	$20 \leq ? \leq 30$	20	$40 \leq ? \leq 105$	70	$140 \leq ? \leq 378$	252

## References

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