# Introduction to Domain Decomposition Methods 

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## MOTIVATION

- DD can be used in the framework of any discretization method for PDEs (FEM, FV, FD, SEM) to make their algebraic solution more efficient on parallel computer platforms
- DDM allow the reformulation of a boundary-value problem on a partition of the computational domain into subdomains $\Rightarrow$ very convenient framework for the solution of heterogeneous or multiphysics problems, i.e. those that are governed by differential equations of different kinds in different subregions of the computational domain.


## THE IDEA

By DDM methods the computational domain $\Omega$ where the bvp is set is subdivided into two or more subdomains on which discretized problems of smaller dimension are to be solved.

Parallel solution algorithms may be used.
There are two ways of subdividing the computational domain:

- with disjoint subdomains
- with overlapping subdomains


Correspondingly, different DD algorithms will be set up.

## EXAMPLES OF SUBDIVISIONS IN APPLICATIONS



Mce X

## CLASSICAL ITERATIVE DD METHODS

## MODEL PROBLEM

Consider the model problem:
find $u: \Omega \rightarrow \mathbb{R}$ s.t.

$$
\begin{cases}L u=f & \text { in } \Omega \\ u=0 & \text { on } \partial \Omega\end{cases}
$$

$L$ is a generic second order elliptic operator.
The weak formulation reads:

$$
\text { find } u \in V=H_{0}^{1}(\Omega): \quad a(u, v)=(f, v) \quad \forall v \in V
$$

where $a(\cdot, \cdot)$ is the bilinear form associated with $L$.

## SCHWARZ METHODS

Consider a decomposition of $\Omega$ with overlap:


The iterative method: given $u_{2}^{(0)}$ on $\Gamma_{1}$, for $k \geq 1$ :

$$
\begin{aligned}
& \text { solve } \begin{cases}L u_{1}^{(k)}=f & \text { in } \Omega_{1} \\
u_{1}^{(k)}=u_{2}^{(k-1)} & \text { on } \Gamma_{1} \\
u_{1}^{(k)}=0 & \text { on } \partial \Omega_{1} \backslash \Gamma_{1}\end{cases} \\
& \text { solve }\left\{\begin{array}{ll}
L u_{2}^{(k)}=f & \text { in } \Omega_{2} \\
u_{2}^{(k)}= \begin{cases}u_{1}^{(k)} & \text { on } \Gamma_{2} \\
u_{1}^{(k-1)} & \text { on } \partial \Omega_{2} \backslash \Gamma_{2} \\
u_{2}^{(k)}=0 & \end{cases}
\end{array} . \begin{array}{ll}
\end{array}\right. \\
& \text { a }
\end{aligned}
$$

Choice of the trace on $\Gamma_{2}$ :

- if $u_{1}^{(k)} \Rightarrow$ multiplicative Schwarz method
- if $u_{1}^{(k-1)} \Rightarrow$ additive Schwarz method

We have two elliptic bvp with Dirichlet conditions in $\Omega_{1}$ and $\Omega_{2}$, and we wish the two sequences $\left\{u_{1}^{(k)}\right\}$ and $\left\{u_{2}^{(k)}\right\}$ to converge to the restrictions of the solution $u$ of the original problem:

$$
\lim _{k \rightarrow \infty} u_{1}^{(k)}=u_{\left.\right|_{1}} \text { and } \quad \lim _{k \rightarrow \infty} u_{2}^{(k)}=u_{\mid \Omega_{2}}
$$

The Schwarz method applied to the model problem always converges, with a rate that increases as the measure $\left|\Gamma_{12}\right|$ of the overlapping region $\Gamma_{12}$ increases.

Example. Consider the model problem

$$
\left\{\begin{array}{l}
-u^{\prime \prime}(x)=0 \\
u(a)=u(b)=0,
\end{array} \quad a<x<b\right.
$$

where


The solution is $u=0$. We show a few iterations of the method:


Clearly, the method converges with a rate that reduces as the length of the interval $\left(\gamma_{2}, \gamma_{1}\right)$ gets smaller.

## NON-OVERLAPPING DECOMPOSITION

We partition now the domain $\Omega$ in two disjoint subdomains:


The following equivalence result holds.

## Theorem

The solution $u$ of the model problem is such that $u_{\left.\right|_{\Omega_{i}}}=u_{i}$ for $i=1,2$, where $u_{i}$ is the solution to the problem

$$
\begin{cases}L u_{i}=f & \text { in } \Omega_{i} \\ u_{i}=0 & \text { on } \partial \Omega_{i} \backslash \Gamma\end{cases}
$$

with interface conditions

$$
u_{1}=u_{2} \quad \text { and } \quad \frac{\partial u_{1}}{\partial n}=\frac{\partial u_{2}}{\partial n} \quad \text { on } \Gamma .
$$

## DIRICHLET-NEUMANN METHOD

Given $u_{2}^{(0)}$ on $\Gamma$, for $k \geq 1$ solve the problems:

$$
\begin{aligned}
& \begin{cases}L u_{1}^{(k)}=f & \text { in } \Omega_{1} \\
u_{1}^{(k)}=u_{2}^{(k-1)} & \text { on } \Gamma \\
u_{1}^{(k)}=0 & \text { on } \partial \Omega_{1} \backslash \Gamma\end{cases} \\
& \begin{cases}L u_{2}^{(k)}=f & \text { in } \Omega_{2} \\
\frac{\partial u_{2}^{(k)}}{\partial n}=\frac{\partial u_{1}^{(k)}}{\partial n} & \text { on } \Gamma \\
u_{2}^{(k)}=0 & \text { on } \partial \Omega_{2} \backslash \Gamma .\end{cases}
\end{aligned}
$$

The equivalence theorem guarantees that when the two sequences $\left\{u_{1}^{(k)}\right\}$ and $\left\{u_{2}^{(k)}\right\}$ converge, then their limit will be necessarily the solution to the exact problem.
The DN algorithm is therefore consistent. Its convergence however is not always guaranteed.

Example. Let $\Omega=(a, b), \gamma \in(a, b), L=-d^{2} / d x^{2}$ and $f=0$. At every $k \geq 1$ the DN algorithm generates the two subproblems:

$$
\begin{gathered}
\begin{cases}-\left(u_{1}^{(k)}\right)^{\prime \prime}=0 & a<x<\gamma \\
u_{1}^{(k)}=u_{2}^{(k-1)} & x=\gamma \\
u_{1}^{(k)}=0 & x=a\end{cases} \\
\begin{cases}-\left(u_{2}^{(k)}\right)^{\prime \prime}=0 & \gamma<x<b \\
\left(u_{2}^{(k)}\right)^{\prime}=\left(u_{1}^{(k)}\right)^{\prime} & x=\gamma \\
u_{2}^{(k)}=0 & x=b .\end{cases}
\end{gathered}
$$

The two sequences converge only if $\gamma>(a+b) / 2$ :


A variant of the DN algorithm can be set up by replacing the Dirichlet condition in the first subdomain by

$$
u_{1}^{(k)}=\theta u_{2}^{(k-1)}+(1-\theta) u_{1}^{(k-1)} \quad \text { on } \Gamma
$$

that is by introducing a relaxation which depends on a positive parameter $\theta$.
In such a way it is always possible to reduce the error between two subsequent iterates.
In the previous example, we can easily verify that, by choosing

$$
\theta_{o p t}=-\frac{u_{1}^{(k-1)}}{u_{2}^{(k-1)}-u_{1}^{(k-1)}},
$$

the algorithm converges to the exact solution in a single iteration.
More in general, there exists a suitable value $\theta_{\max }<1$ such that the DN algorithm converges for any possible choice of the relaxation parameter $\theta$ in the interval $\left(0, \theta_{\text {max }}\right)$.

## NEUMANN-NEUMANN ALGORITHM

Consider again a partition of $\Omega$ into two disjoint subdomains and denote by $\lambda$ the (unknown) value of the solution $u$ at their interface $\Gamma$.

Consider the following iterative algorithm: for any given $\lambda^{(0)}$ on $\Gamma$, for $k \geq 0$ and $i=1,2$, solve the following problems:

$$
\begin{gathered}
\begin{cases}-\triangle u_{i}^{(k+1)}=f & \text { in } \Omega_{i} \\
u_{i}^{(k+1)}=\lambda^{(k)} & \text { on } \Gamma \\
u_{i}^{(k+1)}=0 & \text { on } \partial \Omega_{i} \backslash \Gamma,\end{cases} \\
\begin{cases}-\triangle \psi_{i}^{(k+1)}=0 \\
\frac{\partial \psi_{i}^{(k+1)}}{\partial n}=\frac{\partial u_{1}^{(k+1)}}{\partial n}-\frac{\partial u_{2}^{(k+1)}}{\partial n} & \text { in } \Omega_{i} \\
\psi_{i}^{(k+1)}=0 & \text { on } \Gamma\end{cases} \\
\text { on } \partial \Omega_{i} \backslash \Gamma,
\end{gathered}
$$

with

$$
\lambda^{(k+1)}=\lambda^{(k)}-\theta\left(\sigma_{1} \psi_{1 \mid \Gamma}^{(k+1)}-\sigma_{2} \psi_{2 \mid \Gamma}^{(k+1)}\right)
$$

where $\theta$ is a positive acceleration parameter, while $\sigma_{1}$ and $\sigma_{2}$ are two MC8sitive coefficients.

## ROBIN-ROBIN ALGORITHM

For every $k \geq 0$ solve the following problems:

$$
\begin{cases}-\triangle u_{1}^{(k+1)}=f & \text { in } \Omega_{1} \\ u_{1}^{(k+1)}=0 & \text { on } \partial \Omega_{1} \backslash \Gamma \\ \frac{\partial u_{1}^{(k+1)}}{\partial n}+\gamma_{1} u_{1}^{(k+1)}=\frac{\partial u_{2}^{(k)}}{\partial n}+\gamma_{1} u_{2}^{(k)} & \text { on } \Gamma\end{cases}
$$

then

$$
\begin{cases}-\triangle u_{2}^{(k+1)}=f & \text { in } \Omega_{2} \\ u_{2}^{(k+1)}=0 & \text { on } \partial \Omega_{2} \backslash \Gamma \\ \frac{\partial u_{2}^{(k+1)}}{\partial n}+\gamma_{2} u_{2}^{(k+1)}=\frac{\partial u_{1}^{(k+1)}}{\partial n}+\gamma_{2} u_{1}^{(k+1)} & \text { on } \Gamma\end{cases}
$$

where $u_{2}^{(0)}$ is assigned and $\gamma_{1}, \gamma_{2}$ are non-negative acceleration parameters that satisfy $\gamma_{1}+\gamma_{2}>0$.
Aiming at parallelization, we could use $u_{1}^{(k)}$ instead of $u_{1}^{(k+1)}$.

## THE STEKLOV-POINCARÉ INTERFACE EQUATION

## MULTI-DOMAIN FORMULATION OF POISSON PROBLEM AND INTERFACE CONDITIONS

We consider now the model problem:

$$
\begin{cases}-\triangle u=f & \text { in } \Omega \\ u=0 & \text { on } \partial \Omega\end{cases}
$$

For a domain partitioned into two disjoint subdomains, we can write the equivalent multidomain formulation $\left(u_{i}=u_{\mid \Omega_{i}}, i=1,2\right)$ :

$$
\begin{cases}-\triangle u_{1}=f & \text { in } \Omega_{1} \\ u_{1}=0 & \text { on } \partial \Omega_{1} \backslash \Gamma \\ -\triangle u_{2}=f & \text { in } \Omega_{2} \\ u_{2}=0 & \text { on } \partial \Omega_{2} \backslash \Gamma \\ u_{1}=u_{2} & \text { on } \Gamma \\ \frac{\partial u_{1}}{\partial n}=\frac{\partial u_{2}}{\partial n} & \text { on } \Gamma .\end{cases}
$$

## THE STEKLOV-POINCARÉ OPERATOR

Let $\lambda$ be the unknown value of the solution $u$ on the interface $\Gamma: \lambda=u_{\mid r}$.
Should we know a priori the value $\lambda$ on $\Gamma$, we could solve the following two independent boundary-value problems with Dirichlet condition on $\Gamma$ $(i=1,2)$ :

$$
\begin{cases}-\triangle w_{i}=f & \text { in } \Omega_{i}, \\ w_{i}=0 & \text { on } \partial \Omega_{i} \backslash \Gamma \\ w_{i}=\lambda & \text { on } \Gamma .\end{cases}
$$

With the aim of obtaining the value $\lambda$ on $\Gamma$, let us split $w_{i}$ as follows

$$
w_{i}=w_{i}^{*}+u_{i}^{0}
$$

where $w_{i}^{*}$ and $u_{i}^{0}$ represent the solutions of the following problems ( $i=1,2$ ):

$$
\begin{cases}-\triangle w_{i}^{*}=f & \text { in } \Omega_{i}, \\ w_{i}^{*}=0 & \text { on } \partial \Omega_{i} \cap \partial \Omega \\ w_{i}^{*}=0 & \text { on } \Gamma\end{cases}
$$

and

$$
\begin{cases}-\triangle u_{i}^{0}=0 & \text { in } \Omega_{i} \\ u_{i}^{0}=0 & \text { on } \partial \Omega_{i} \cap \partial \Omega \\ u_{i}^{0}=\lambda & \text { on } \Gamma\end{cases}
$$

- The functions $w_{i}^{*}$ depend solely on the source data $f \Rightarrow w_{i}^{*}=G_{i} f$ where $G_{i}$ is a linear continuous operator
- $u_{i}^{0}$ depend solely on the value $\lambda$ on $\Gamma \Rightarrow u_{i}^{0}=H_{i} \lambda$, where $H_{i}$ is the so-called harmonic extension operator of $\lambda$ on the domain $\Omega_{i}$.
We have that

$$
u_{i}=w_{i}^{*}+u_{i}^{0},(i=1,2,) \quad \Leftrightarrow \quad \frac{\partial w_{1}}{\partial n}=\frac{\partial w_{2}}{\partial n} \quad \text { on } \Gamma \text {. }
$$

Using the previously introduced notations:

$$
\frac{\partial}{\partial n}\left(G_{1} f+H_{1} \lambda\right)=\frac{\partial}{\partial n}\left(G_{2} f+H_{2} \lambda\right),
$$

and therefore

$$
\left(\frac{\partial H_{1}}{\partial n}-\frac{\partial H_{2}}{\partial n}\right) \lambda=\left(\frac{\partial G_{2}}{\partial n}-\frac{\partial G_{1}}{\partial n}\right) f \quad \text { on } \Gamma .
$$

We have obtained the Steklov-Poincaré equation for the unknown $\lambda$ on the interface $\Gamma$ :

$$
S \lambda=\chi \quad \text { on } \Gamma
$$

- $S$ is the Steklov-Poincaré pseudo-differential operator:

$$
S_{\mu}=\frac{\partial}{\partial n} H_{1} \mu-\frac{\partial}{\partial n} H_{2} \mu=\sum_{i=1}^{2} \frac{\partial}{\partial n_{i}} H_{i} \mu
$$

- $\chi$ is a linear functional which depends on $f$ :

$$
\chi=\frac{\partial}{\partial n} G_{2} f-\frac{\partial}{\partial n} G_{1} f=-\sum_{i=1}^{2} \frac{\partial}{\partial n_{i}} G_{i} f .
$$

- The operator

$$
S_{i}: \mu \rightarrow S_{i} \mu=\left.\frac{\partial}{\partial n_{i}}\left(H_{i} \mu\right)\right|_{\Gamma}, \quad i=1,2
$$

is called local Steklov-Poincaré operator which operates between the trace space

$$
\Lambda=\left\{\mu: \exists v \in V \text { s.t. } \mu=v_{\mid \Gamma}\right\}=H_{00}^{1 / 2}(\Gamma)
$$

and its dual $\Lambda^{\prime}$.

## EQUIVALENCE BETWEEN THE DD SCHEMES AND CLASSICAL ITERATIVE METHODS

- The DN method can be reinterpreted as a preconditioned Richardson method for the solution of the Steklov-Poincaré interface equation:

$$
P_{D N}\left(\lambda^{(k)}-\lambda^{(k-1)}\right)=\theta\left(\chi-S \lambda^{(k-1)}\right)
$$

The preconditioning operator is $P_{D N}=S_{2}=\partial\left(H_{2} \mu\right) / \partial n_{2}$.

- The NN method can also be interpreted as a preconditioned Richardson algorithm

$$
P_{N N}\left(\lambda^{(k)}-\lambda^{(k-1)}\right)=\theta\left(\chi-S \lambda^{(k-1)}\right)
$$

with $P_{N N}=\left(\sigma_{1} S_{1}^{-1}+\sigma_{2} S_{2}^{-1}\right)^{-1}$.

- The Robin-Robin algorithm is equivalent to the following alternating direction (ADI) algorithm:

$$
\begin{aligned}
\left(\gamma_{1} \mathcal{I}_{\Lambda}+S_{1}\right) \mu_{1}^{(k)} & =\chi+\left(\gamma_{1} \mathcal{I}_{\Lambda}+S_{2}\right) \mu_{2}^{(k-1)} \\
\left(\gamma_{2} \mathcal{I}_{\Lambda}+S_{2}\right) \mu_{2}^{(k)} & =\chi+\left(\gamma_{2} \mathcal{I}_{\Lambda}+S_{1}\right) \mu_{1}^{(k-1)}
\end{aligned}
$$

where $\mathcal{I}_{\Lambda}: \Lambda \rightarrow \Lambda^{\prime}$ here denotes the Riesz isomorphism between the Hilbert space $\Lambda$ and its dual $\Lambda^{\prime}$.

At convergence (for a convenient choice of $\gamma_{1}$ and $\gamma_{2}$ ), we have $\mu_{1}=\mu_{2}=\lambda$.

## FINITE ELEMENT APPROXIMATION: MULTIDOMAIN FORMULATION

The Galerkin finite element approximation of the Poisson problem on a triangulation $\mathcal{T}_{h}$ of $\Omega$ reads

$$
\text { find } u_{h} \in V_{h}: \quad a\left(u_{h}, v_{h}\right)=F\left(v_{h}\right) \quad \forall v_{h} \in V_{h},
$$

where

$$
V_{h}=\left\{v_{h} \in C^{0}(\bar{\Omega}): v_{\left.h\right|_{K}} \in \mathbb{P}_{r} \quad r \geq 1, \forall K \in \mathcal{T}_{h}, v_{h}=0 \text { on } \partial \Omega\right\}
$$

is the space of finite element functions of degree $r$ with basis $\left\{\varphi_{j}\right\}_{j=1}^{N_{h}}$.
We partition the nodes of the triangulation as follows:

- $\left\{x_{j}^{(1)}, 1 \leq j \leq N_{1}\right\}$ nodes in subdomain $\Omega_{1}$
- $\left\{x_{j}^{(2)}, 1 \leq j \leq N_{2}\right\}$ nodes in subdomain $\Omega_{2}$
- $\left\{x_{j}^{(\Gamma)}, 1 \leq j \leq N_{\Gamma}\right\}$ nodes on on the interface $\Gamma$


## Example:



We split the basis functions accordingly:

- $\varphi_{j}^{(1)}$ functions associated to the nodes $x_{j}^{(1)}$
- $\varphi_{j}^{(2)}$ functions associated to the nodes $x_{j}^{(2)}$
- $\varphi_{j}^{(\Gamma)}$ functions associated to the nodes $x_{j}^{(\Gamma)}$ on the interface

Then, we can write:

$$
\begin{aligned}
u_{h}(x)= & \sum_{j=1}^{N_{1}} u_{h}\left(x_{j}^{(1)}\right) \varphi_{j}^{(1)}(x)+\sum_{j=1}^{N_{2}} u_{h}\left(x_{j}^{(2)}\right) \varphi_{j}^{(2)}(x) \\
& +\sum_{j=1}^{N_{\Gamma}} u_{h}\left(x_{j}^{(\Gamma)}\right) \varphi_{j}^{(\Gamma)}(x)
\end{aligned}
$$

Substituing in the Galerkin formulation, after some algebra, we end up with the linear system

$$
A \mathbf{u}=\mathbf{f}, \text { that is }\left[\begin{array}{ccc}
A_{11} & 0 & A_{1 \Gamma} \\
0 & A_{22} & A_{2 \Gamma} \\
A_{\Gamma 1} & A_{\Gamma 2} & A_{\Gamma \Gamma}
\end{array}\right]\left[\begin{array}{c}
\mathbf{u}_{1} \\
\mathbf{u}_{2} \\
\boldsymbol{\lambda}
\end{array}\right]=\left[\begin{array}{c}
\mathbf{f}_{1} \\
\mathbf{f}_{2} \\
\mathbf{f}_{\Gamma}
\end{array}\right]
$$

where $A_{\Gamma \Gamma}=\left(A_{\Gamma \Gamma}^{(1)}+A_{\Gamma \Gamma}^{(2)}\right)$ and $\mathbf{f}_{\Gamma}=\mathbf{f}_{1}^{\Gamma}+\mathbf{f}_{2}^{\Gamma}$.
Moreover,

$$
\mathbf{u}_{1}=\left(u_{h}\left(x_{j}^{(1)}\right)\right), \mathbf{u}_{2}=\left(u_{h}\left(x_{j}^{(2)}\right)\right) \text { and } \boldsymbol{\lambda}=\left(u_{h}\left(x_{j}^{(\Gamma)}\right)\right)
$$

## THE SCHUR COMPLEMENT SYSTEM

Since $\lambda$ represents the unknown value of $u$ on $\Gamma$, its finite element correspondent is the vector $\boldsymbol{\lambda}$ of the values of $u_{h}$ at the interface nodes.

By Gaussian elimination, we can obtain a new reduced system on the sole unknown $\boldsymbol{\lambda}$ :

- Matrices $A_{11}$ and $A_{22}$ are invertible since they are associated with two homogeneous Dirichlet boundary-value problems for the Laplace operator:

$$
\mathbf{u}_{1}=A_{11}^{-1}\left(\mathbf{f}_{1}-A_{1 \Gamma} \boldsymbol{\lambda}\right) \quad \text { and } \quad \mathbf{u}_{2}=A_{22}^{-1}\left(\mathbf{f}_{2}-A_{2 \Gamma} \boldsymbol{\lambda}\right)
$$

- From the third equation we obtain:

$$
\begin{aligned}
& {\left[\left(A_{\Gamma \Gamma}^{(1)}-A_{\Gamma 1} A_{11}^{-1} A_{1 \Gamma}\right)+\left(A_{\Gamma \Gamma}^{(2)}-A_{\Gamma 2} A_{22}^{-1} A_{2 \Gamma}\right)\right] \boldsymbol{\lambda}=} \\
& \mathbf{f}_{\Gamma}-A_{\Gamma 1} A_{11}^{-1} \mathbf{f}_{1}-A_{\Gamma 2} A_{22}^{-1} \mathbf{f}_{2} .
\end{aligned}
$$

Setting

$$
\Sigma=\Sigma_{1}+\Sigma_{2} \quad \text { with } \quad \Sigma_{i}=A_{\Gamma \Gamma}^{(i)}-A_{\Gamma i} A_{i i}^{-1} A_{i \Gamma} \quad(i=1,2)
$$

and

$$
\chi_{\Gamma}=\mathbf{f}_{\Gamma}-A_{\Gamma 1} A_{11}^{-1} \mathbf{f}_{1}-A_{\Gamma 2} A_{22}^{-1} \mathbf{f}_{2}
$$

we obtain the Schur complement system

$$
\Sigma \boldsymbol{\lambda}=\chi_{\Gamma}
$$

- $\Sigma$ and $\chi_{\Gamma}$ approximate $S$ and $\chi$.
- $\Sigma$ is the so-called Schur complement of $A$ with respect to $\mathbf{u}_{1}$ and $\mathbf{u}_{2}$.
- $\Sigma_{i}$ are the Schur complements related to the subdomains $\Omega_{i}$ ( $i=1,2$ ).


## Example:



On the left, decomposition of $\Omega=(0,1) \times(0,1)$ into four square subdomains. On the right, sparsity pattern of the Schur complement matrix arising from the decomposition depicted on the left.

## PROPERTIES OF THE SCHUR COMPLEMENT $\Sigma$

The Schur complement $\Sigma$ inherits some of the properties of $A$ :

- if $A$ is singular, so is $\Sigma$;
- if $A$ (respectively, $A_{i i}$ ) is symmetric, then $\Sigma$ (respectively, $\Sigma_{i}$ ) is symmetric too;
- if $A$ is positive definite, so is $\Sigma$.

Moreover, concerning the condition number, we have

- $K_{2}(A) \simeq C h^{-2}$
- $K_{2}(\Sigma) \simeq C h^{-1}$


## EQUIVALENCE BETWEEN THE DN METHOD AND PRECONDITIONER RICHARDSON IN THE DISCRETE CASE

In the discrete case it is easy to prove that the DN method is equivalent to a preconditioned Richardson algorithm.

First, we write the algebraic DN method:

- Dirichlet problem:

$$
A_{11} \mathbf{u}_{1}^{(k)}=\mathbf{f}_{1}-A_{1 \Gamma} \boldsymbol{\lambda}^{(k-1)}
$$

- Neumann problem

$$
\left[\begin{array}{ll}
A_{22} & A_{2 \Gamma} \\
A_{\Gamma 2} & A_{\Gamma \Gamma}^{(2)}
\end{array}\right]\left[\begin{array}{c}
\mathbf{u}_{2}^{(k)} \\
\boldsymbol{\lambda}^{(k-1 / 2)}
\end{array}\right]=\left[\begin{array}{c}
\mathbf{f}_{2} \\
\mathbf{f}_{\Gamma}-A_{\Gamma 1} \mathbf{u}_{1}^{(k)}-A_{\Gamma \Gamma}^{(1)} \boldsymbol{\lambda}^{(k-1)}
\end{array}\right]
$$

- Relaxation:

$$
\boldsymbol{\lambda}^{(k)}=\theta \boldsymbol{\lambda}^{(k-1 / 2)}+(1-\theta) \boldsymbol{\lambda}^{(k-1)}
$$

By eliminating $\mathbf{u}_{2}^{(k)}$ we obtain:

$$
\left(A_{\Gamma \Gamma}^{(2)}-A_{\Gamma 2} A_{22}^{-1} A_{2 \Gamma}\right) \boldsymbol{\lambda}^{(k-1 / 2)}=\mathbf{f}_{\Gamma}-A_{\Gamma 1} \mathbf{u}_{1}^{(k)}-A_{\Gamma \Gamma}^{(1)} \boldsymbol{\lambda}^{(k-1)}-A_{\Gamma 2} A_{22}^{-1} \mathbf{f}_{2}
$$

and using the definition of $\Sigma_{2}$ we have

$$
\Sigma_{2} \boldsymbol{\lambda}^{(k-1 / 2)}=\mathbf{f}_{\Gamma}-A_{\Gamma 1} A_{11}^{-1} \mathbf{f}_{1}-A_{\Gamma 2} A_{22}^{-1} \mathbf{f}_{2}-\left(A_{\Gamma \Gamma}^{(1)}-A_{\Gamma 1} A_{11}^{-1} A_{1 \Gamma}\right) \boldsymbol{\lambda}^{(k-1)},
$$

that is, owing to the definition of $\Sigma_{1}$ and $\chi_{\Gamma}$ :

$$
\boldsymbol{\lambda}^{(k-1 / 2)}=\Sigma_{2}^{-1}\left(\chi_{\Gamma}-\Sigma_{1} \boldsymbol{\lambda}^{(k-1)}\right)
$$

Then,

$$
\begin{aligned}
\boldsymbol{\lambda}^{(k)} & =\theta \Sigma_{2}^{-1}\left(\chi_{\Gamma}-\Sigma_{1} \boldsymbol{\lambda}^{(k-1)}\right)+(1-\theta) \boldsymbol{\lambda}^{(k-1)} \\
& =\theta \Sigma_{2}^{-1}\left(\chi_{\Gamma}-\Sigma \boldsymbol{\lambda}^{(k-1)}+\Sigma_{2} \boldsymbol{\lambda}^{(k-1)}\right)+(1-\theta) \boldsymbol{\lambda}^{(k-1)}
\end{aligned}
$$

whence

$$
\Sigma_{2}\left(\boldsymbol{\lambda}^{(k)}-\boldsymbol{\lambda}^{(k-1)}\right)=\theta\left(\chi_{\Gamma}-\Sigma \boldsymbol{\lambda}^{(k-1)}\right)
$$

Similarly, we can prove that the other iterative methods that we have illustrated are equivalent to preconditioned Richiardson methods for the Schur complement system with preconditioners:

- for the DN algorithm: $P_{h}=\Sigma_{2}$
- for the ND algorithm: $P_{h}=\Sigma_{1}$
- for the NN algorithm: $P_{h}=\left(\sigma_{1} \Sigma_{1}^{-1}+\sigma_{2} \Sigma_{2}^{-1}\right)^{-1}$
- for the RR algorithm: $P_{h}=\left(\gamma_{1}+\gamma_{2}\right)^{-1}\left(\gamma_{1} I+\Sigma_{1}\right)\left(\gamma_{2} I+\Sigma_{2}\right)$

All these preconditioners are optimal.

This is not true in the case of more subdomains, where the condition number of the Schur complement matrix $\Sigma$ is

$$
\kappa_{2}(\Sigma) \leq C \frac{H}{h H_{\min }^{2}}
$$

where $H$ and $H_{\text {min }}$ are the maximal and minimal diameters of the subdomains, respectively.

| $\kappa(\Sigma)$ | $H=1 / 2$ | $H=1 / 4$ | $H=1 / 8$ |
| :---: | :---: | :---: | :---: |
| $\mathrm{~h}=1 / 8$ | 9.77 | 14.83 | 25.27 |
| $\mathrm{~h}=1 / 16$ | 21.49 | 35.25 | 58.60 |
| $\mathrm{~h}=1 / 32$ | 44.09 | 75.10 | 137.73 |
| $\mathrm{~h}=1 / 64$ | 91.98 | 155.19 | 290.43 |

Condition number of the Schur complement matrix.

NON-OVERLAPPING METHODS IN THE CASE OF MORE THAN TWO SUBDOMAINS

## SCALABILITY

Definition A preconditioner $P_{h}$ of $\Sigma$ is said to be scalable if the condition number of the preconditioned matrix $P_{h}^{-1} \Sigma$ is independent of the number of subdomains.

Iterative methods using scalable preconditioners allow henceforth to achieve convergence rates independent of the subdomain number.

## RESTRICTION AND EXTENSION OPERATORS

We introduce a restriction operator $R_{i}$ that, to any vector $\mathbf{v}_{h}$ of nodal values on the global domain $\Omega$, associates its restriction to the subdomain $\Omega_{i}$ :

$$
R_{i}: \mathbf{v}_{\left.h\right|_{\Omega}} \rightarrow \mathbf{v}_{h \mid \Omega_{i} u r_{i}}^{i} .
$$

Moreover $R_{i}^{T}$

$$
R_{i}^{T}: \mathbf{v}_{h \Omega_{\Omega_{i}} u_{r}}^{i} \rightarrow \mathbf{v}_{\left.h\right|_{\Omega}}
$$

is the prolongation (or extension-by-zero) operator.
In algebraic form $R_{i}$ can be represented by a matrix that coincides with the identity matrix in correspondence with the subdomain $\Omega_{i}$ :

$$
R_{i}=\left[\begin{array}{ccc|ccc|ccc}
0 & \ldots & 0 & & 1 & 0 & \ldots & 0 \\
\vdots & \ddots & \vdots & & \ddots & & \vdots & \ddots & \vdots \\
0 & \ldots & 0 & & \underbrace{}_{\Omega_{i}} & 0 & \ldots & 0
\end{array}\right]
$$

Similarly we can define the restriction and prolongation operators $R_{\Gamma_{i}}$ and $R_{\Gamma_{i}}^{T}$ that act on the vector of the interface nodal values.

A possible preconditioner for $\Sigma$ is

$$
P_{h}=\sum_{i=1}^{M} R_{\Gamma_{i}}^{T} \Sigma_{i} R_{\Gamma_{i}}
$$

More in general, we combine contributions of local subdmain preconditioners with that of a global contribution referring to a coarse grid whose elements are the subdomains themselves:

$$
\left(P_{h}\right)^{-1}=\sum_{i=1}^{M} R_{\Gamma_{i}}^{T} P_{i, h}^{-1} R_{\Gamma_{i}}+R_{\Gamma}^{T} P_{H}^{-1} R_{\Gamma}
$$

## NEUMANN-NEUMANN PRECONDITIONER

The Neumann-Neumann preconditioner for more subdomains reads:

$$
\left(P_{h}^{N N}\right)^{-1}=\sum_{i=1}^{M} R_{\Gamma_{i}}^{T} D_{i} \Sigma_{i}^{*} D_{i} R_{\Gamma_{i}}
$$

where $\Sigma_{i}^{*}$ is either $\Sigma_{i}^{-1}$ or an approximation of $\Sigma_{i}^{-1}$.
$D_{i}$ is a diagonal matrix of positive weights

$$
D_{i}=\left[\begin{array}{lll}
d_{1} & & \\
& \ddots & \\
& & d_{n}
\end{array}\right]
$$

$d_{j}$ is the number of subdomains that share the $j$-th node.
We have the following estimate:

$$
\kappa\left(\left(P_{h}^{N N}\right)^{-1} \Sigma\right) \leq C H^{-2}\left(1+\log \frac{H}{h}\right)^{2}
$$

## Example

| $\kappa\left(\left(P_{h}^{N N}\right)^{-1} \Sigma\right)$ | $H=1 / 2$ | $H=1 / 4$ | $H=1 / 8$ | $H=1 / 16$ |
| :---: | :---: | :---: | :---: | :---: |
| $h=1 / 16$ | 2.55 | 15.20 | 47.60 | - |
| $h=1 / 32$ | 3.45 | 20.67 | 76.46 | 194.65 |
| $h=1 / 64$ | 4.53 | 26.25 | 105.38 | 316.54 |
| $h=1 / 128$ | 5.79 | 31.95 | 134.02 | 438.02 |
| Condition number of $\left(P_{h}^{N N}\right)^{-1} \Sigma$ |  |  |  |  |

## BALANCED NEUMANN-NEUMANN PRECONDITIONER

- The Neumann-Neumann preconditioner of the Schur complement system is not scalable.
- A substantial improvement can be achieved by adding a coarse grid correction:

$$
\left(P_{h}^{B N N}\right)^{-1}=\Sigma_{H}^{-1}+\left(I-\Sigma_{H}^{-1} \Sigma\right)\left(P_{h}^{N N}\right)^{-1}\left(I-\Sigma \Sigma_{H}^{-1}\right)
$$

where $\Sigma_{H}^{-1}=R_{\Gamma}^{T} A_{H}^{-1} R_{\Gamma}$.
This is called balanced Neumann-Neumann preconditioner.

- We can prove that

$$
\kappa_{2}\left(\left(P_{h}^{B N N}\right)^{-1} \Sigma\right) \leq C\left(1+\log \frac{H}{h}\right)^{2}
$$

## Example:

| $\kappa\left(\left(P_{h}^{B N N}\right)^{-1} \Sigma\right)$ | $H=1 / 2$ | $H=1 / 4$ | $H=1 / 8$ | $H=1 / 16$ |
| :---: | :---: | :---: | :---: | :---: |
| $h=1 / 16$ | 1.67 | 1.48 | 1.27 | - |
| $h=1 / 32$ | 2.17 | 2.03 | 1.47 | 1.29 |
| $h=1 / 64$ | 2.78 | 2.76 | 2.08 | 1.55 |
| $h=1 / 128$ | 3.51 | 3.67 | 2.81 | 2.07 |

## ALGEBRAIC FORM OF SCHWARZ ITERATIVE METHODS

## ALGEBRAIC FORM OF SCHWARZ METHODS FOR FE DISCRETIZATION

- Let $A \mathbf{u}=\mathbf{f}$ be the system associated to the finite element approximation of the Poisson problem.
- We still assume that $\Omega$ is decomposed in two overlapping subdomains $\Omega_{1}$ and $\Omega_{2}$.
- We denote by $N_{h}$ the total number of interior nodes of $\Omega$, and by $n_{1}$ and $n_{2}$ the interior nodes of $\Omega_{1}$ and $\Omega_{2}: N_{h} \leq n_{1}+n_{2}$
- The stiffness matrix $A$ contains two submatrices, say $A_{1}$ and $A_{2}$, which correspond to the local stiffness matrices associated to the Dirichlet problems in $\Omega_{1}$ and $\Omega_{2}$ :

- We have:

$$
A_{i}=R_{i} A R_{i}^{T} \in \mathbb{R}^{n_{i} \times n_{i}}
$$

with

$$
R_{1}^{T}=\left[\begin{array}{ccc}
1 & \ldots & 0 \\
\vdots & \ddots & \vdots \\
0 & \ldots & 1 \\
& \mathbf{0} & \\
& &
\end{array}\right] \in \mathbb{R}^{N_{h} \times n_{1}} \quad \text { and } R_{2}^{T}=\left[\begin{array}{ccc} 
& & \\
& \mathbf{0} & \\
1 & \ldots & 0 \\
\vdots & \ddots & \vdots \\
0 & \ldots & 1
\end{array}\right] \in \mathbb{R}^{N_{h} \times n_{2}} .
$$

## THE MULTIPLICATIVE SCHWARZ METHOD

Using these definitions, one iteration of the multiplicative Schwarz method for the system $A \mathbf{u}=\mathbf{f}$ can be represented as:

$$
\begin{aligned}
& \mathbf{u}^{k+1 / 2}=\mathbf{u}^{k}+R_{1}^{T} A_{1}^{-1} R_{1}\left(\mathbf{f}-A \mathbf{u}^{k}\right) \\
& \mathbf{u}^{k+1}=\mathbf{u}^{k+1 / 2}+R_{2}^{T} A_{2}^{-1} R_{2}\left(\mathbf{f}-A \mathbf{u}^{k+1 / 2}\right)
\end{aligned}
$$

Setting

$$
P_{i}=R_{i}^{T} A_{i}^{-1} R_{i} A
$$

this corresponds to

$$
\begin{aligned}
\mathbf{u}^{k+1 / 2} & =\left(I-P_{1}\right) \mathbf{u}^{k}+P_{1} \mathbf{u} \\
\mathbf{u}^{k+1}=\left(I-P_{2}\right) \mathbf{u}^{k+1 / 2}+P_{2} \mathbf{u} & =\left(I-P_{2}\right)\left(I-P_{1}\right) \mathbf{u}^{k}+\left(P_{1}+P_{2}-P_{2} P_{1}\right) \mathbf{u}
\end{aligned}
$$

If we define, for $i=1,2$

$$
Q_{i}:=R_{i}^{T} A_{i}^{-1} R_{i}=P_{i} A^{-1}
$$

then we obtain

$$
\begin{aligned}
\mathbf{u}^{k+1} & =\mathbf{u}^{k}+Q_{1}\left(\mathbf{f}-A \mathbf{u}^{k}\right)+Q_{2}\left[\mathbf{f}-A\left(\mathbf{u}^{k}+Q_{1}\left(\mathbf{f}-A \mathbf{u}^{k}\right)\right)\right] \\
& =\mathbf{u}^{k}+\left(Q_{1}+Q_{2}-Q_{2} Q_{1}\right)\left(\mathbf{f}-A \mathbf{u}^{k}\right)
\end{aligned}
$$

## THE ADDITIVE SCHWARZ METHOD

One iteration of the additive Schwarz method becomes:

$$
\mathbf{u}^{(k+1)}=\mathbf{u}^{(k)}+\left(R_{1}^{T} A_{1}^{-1} R_{1}+R_{2}^{T} A_{2}^{-1} R_{2}\right)\left(\mathbf{f}-A \mathbf{u}^{(k)}\right)
$$

and therefore

$$
\mathbf{u}^{(k+1)}=\left(I-P_{1}-P_{2}\right) \mathbf{u}^{(k)}+\left(P_{1}+P_{2}\right) \mathbf{u}
$$

Using again

$$
Q_{i}:=R_{i}^{T} A_{i}^{-1} R_{i}=P_{i} A^{-1}
$$

we obtain

$$
\mathbf{u}^{k+1}=\mathbf{u}^{k}+\left(Q_{1}+Q_{2}\right)\left(\mathbf{f}-A \mathbf{u}^{k}\right)
$$

In the case of a decomposition of $\Omega$ into $M \geq 2$ overlapping subdomains $\left\{\Omega_{i}\right\}$ we have:

$$
\mathbf{u}^{k+1}=\mathbf{u}^{k}+\left(\sum_{i=1}^{M} Q_{i}\right)\left(\mathbf{f}-A \mathbf{u}^{k}\right)
$$

## THE SCHWARZ METHOD AS A PRECONDITIONER

- By defining

$$
P_{\mathrm{as}}=\left(\sum_{i=1}^{M} Q_{i}\right)^{-1}
$$

it follows that one iteration of the additive Schwarz method corresponds to a preconditioned Richardson iteration for system $A \mathbf{u}=\mathbf{f}$ with preconditioner $P_{\text {as }}$ (additive Schwarz preconditioner).

- $P_{a s}$ is not optimal as the condition number blows up if the size of the subdomains reduces:

$$
\kappa_{2}\left(P_{a s}^{-1} A\right) \leq C \frac{1}{\delta H}
$$

where $C$ is a constant independent of $h, H, \delta$.

## Example

| $\kappa\left(P_{a s}^{-1} A\right)$ | $H=1 / 2$ | $H=1 / 4$ | $H=1 / 8$ | $H=1 / 16$ |
| :---: | :---: | :---: | :---: | :---: |
| $h=1 / 16$ | 15.95 | 27.09 | 52.08 | - |
| $h=1 / 32$ | 31.69 | 54.52 | 104.85 | 207.67 |
| $h=1 / 64$ | 63.98 | 109.22 | 210.07 | 416.09 |
| $h=1 / 128$ | 127.99 | 218.48 | 420.04 | 832.57 |

This is due to the fact that the exchange of information occurs only between neighboring subdomains, since only local solves are involved by the application of $\left(P_{a s}\right)^{-1}$.
$\longrightarrow$ We have to introduce a "coarse" global problem over the whole domain to guarantee a mechanism of global communication among all subdomains.

## TWO-LEVEL SCHWARZ PRECONDITIONERS

- As for the Neumann-Neumann method, we can introduce a coarse-grid mechanism that allows for a sudden information diffusion on the whole domain $\Omega$ :
- consider the subdomains as macro-elements of a new coarse grid $\mathcal{T}_{H}$
- build a corresponding stiffness matrix $A_{H}$
- The matrix

$$
Q_{0}=R_{H}^{T} A_{H}^{-1} R_{H}
$$

represents the coarse level correction for the two-level preconditioner, with $R_{H}$ the restriction operator from the fine to the coarse grid.

- The two-level preconditioner $P_{c a s}$ is defined as:

$$
P_{\text {cas }}^{-1}=\sum_{i=0}^{M} Q_{i}
$$

We can prove that there exists a constant $C>0$, independent of both $h$ and $H$ such that

$$
\kappa_{2}\left(P_{\text {cas }}^{-1} A\right) \leq C \frac{H}{\delta}
$$

If $\delta$ is a fraction of $H$, the preconditioner $P_{\text {cas }}$ is scalable.

- Iterations on the original finite element system using $P_{\text {cas }}$ converges with a rate independent of $h$ and $H$ (and therefore of the number of subdomains)
- Thanks to the additive structure of the preconditioner, the preconditioning step is fully parallel as it involves the solution of independent systems, one per each local matrix $A_{i}$.


## Example

| $\kappa\left(P_{\text {cas }}^{-1} A\right)$ | $H=1 / 4$ | $H=1 / 8$ | $H=1 / 16$ |
| :---: | :---: | :---: | :---: |
| $h=1 / 32$ | 7.03 | 4.94 | - |
| $h=1 / 64$ | 12.73 | 7.59 | 4.98 |
| $h=1 / 128$ | 23.62 | 13.17 | 7.66 |
| $h=1 / 256$ | 45.33 | 24.34 | 13.28 |

If $H / \delta=$ constant, this two-level preconditioner is either optimal and scalable.

## Practical indications:

- for decompositions with a small number of subdomains, the single level Schwarz preconditioner $P_{a s}$ is very efficient;
- when the number $M$ of subdomains gets large, using two-level preconditioners becomes crucial.
In those cases in which the generation of a coarse grid is difficult, other algebraic techniques, like aggregation, can be adopted:

$$
\kappa_{2}\left(P_{\text {aggre }}^{-1} A\right) \leq C\left(1+\frac{H}{\delta}\right)
$$

| $P_{\text {aggre }}^{-1} A$ | $H=1 / 4$ | $H=1 / 8$ | $H=1 / 16$ |
| :---: | :---: | :---: | :---: |
| $h=1 / 16$ | 13.37 | 8.87 | - |
| $h=1 / 32$ | 26.93 | 17.71 | 9.82 |
| $h=1 / 64$ | 54.33 | 35.21 | 19.70 |
| $h=1 / 128$ | 109.39 | 70.22 | 39.07 |

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